



Merge Sort


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History of Merge Sort

- Merge sort was one of the first sorting algorithms proposed for computing, brought forward by John von Neumann in 1945.
- Merge sort becomes one of the first divide-and-conquer sorting algorithms
- The algorithm has roots in card-sorting machines of the late 19th century



Von Neumann



The Merge Sort Algorithm

4	3	6	2	7	1	8	5
---	---	---	---	---	---	---	---

```
procedure: mergesort ( $a_1, a_2, \dots, a_n$ : array of integers)
if (length of array  $\leq 1$ ), then
    return array
left := left half of input array
right := right half of input array
left := mergesort(left)
right := mergesort(right)
```

```
final := empty array
while (left is not empty and right is not empty) then
    if (first element of left  $\leq$  first element of right), then
        append first element of left to final
        remove first element of left
    else if (first element of right  $<$  first element of left), then
        append first element of right to final
        remove first element of right
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```
if (left is still not empty), then
    append rest of left to final
else if (right is still not empty), then
    append rest of right to final
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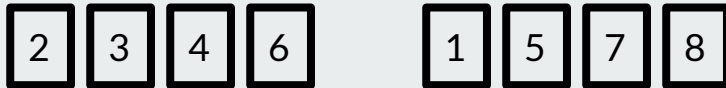
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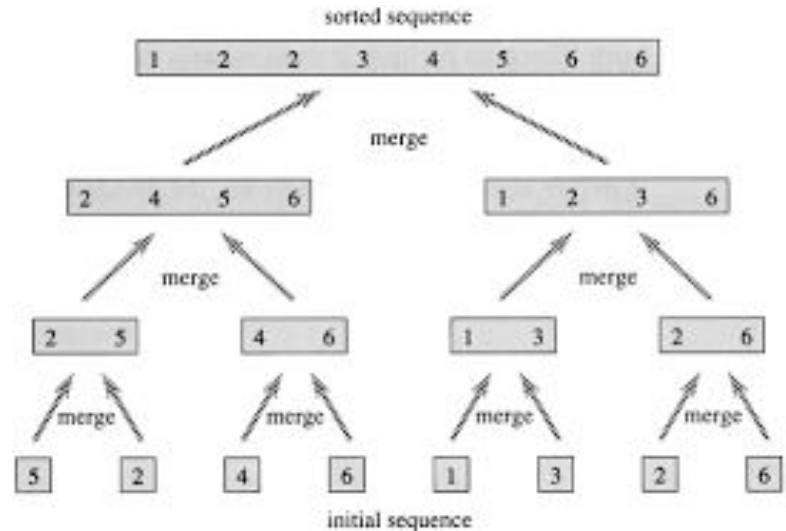


Complexity of Merge sort

Algorithm	Merge Sort	Insertion Sort (a simple sorting algorithm)
Best-case time complexity	$\Omega(n \log n)$	$O(n^2)$
Worst-case time complexity	$O(n \log n)$	$O(n)$ comparisons, $O(1)$ swaps
Space complexity	$O(n)$	$O(n)$

Alternative Merge Sort: Iterative Version

- We just examined a **top-down** implementation of Merge Sort, which uses recursion to sort the entire list
- However, we can implement a **bottom-up** Merge Sort, **without using recursion**
- To do this, we immediately start merging and sorting instead of first dividing the sequence into groups





Bottom-up Merge Sort

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procedure: mergesort ( $a_1, a_2, \dots, a_n$ : array of integers)
list := input array
final := empty array
for (k=0, k<log2(length of list), k++), do
  for (j=0, j<n, j+=2(k+1)), do
    merge(list, final, j, 2k)
```



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  for ( $j=0, j < n, j += 2^{(k+1)}$ ), do
    merge(list, final,  $j, 2^k$ )
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